

SYMMETRIC 1-DEPENDENT COLORINGS OF THE INTEGERS

ALEXANDER E. HOLROYD AND THOMAS M. LIGGETT

ABSTRACT. In a recent paper by the same authors, we constructed a stationary 1-dependent 4-coloring of the integers that is invariant under permutations of the colors. This was the first stationary k -dependent q -coloring for any k and q . When the analogous construction is carried out for $q > 4$ colors, the resulting process is not k -dependent for any k . We construct here a process that is symmetric in the colors and 1-dependent for every $q \geq 4$. The construction uses a recursion involving Chebyshev polynomials evaluated at $\sqrt{q}/2$.

1. INTRODUCTION

By a (proper) q -coloring of the integers, we mean a sequence $(X_i : i \in \mathbb{Z})$ of $[q]$ -valued random variables satisfying $X_i \neq X_{i+1}$ for all i (where $[q] := \{1, \dots, q\}$). The coloring is said to be stationary if the (joint) distribution of $(X_i : i \in \mathbb{Z})$ agrees with that of $(X_{i+1} : i \in \mathbb{Z})$, and k -dependent if the families $(X_i : i \leq m)$ and $(X_i : i > m + k)$ are independent of each other for each m . In [2], we gave a construction of a stationary 1-dependent 4-coloring of the integers that is invariant under permutations of the colors. When the same construction is carried out for $q > 4$ colors, the resulting distribution is not k -dependent for any k . Of course, the 1-dependent 4-coloring is also a 1-dependent q -coloring for every $q > 4$, and one may obtain other 1-dependent q -colorings by splitting a color into further colors using an independent source of randomness. However, these colorings are not symmetric in the colors. We give here a modification of the process of [2] that is symmetric in the colors and 1-dependent for every $q \geq 4$. Here is our main result.

Theorem 1. *For each integer $q \geq 4$, there exists a stationary 1-dependent q -coloring of the integers that is invariant in law under permutations of the colors and under the reflection $(X_i : i \in \mathbb{Z}) \mapsto (X_{-i} : i \in \mathbb{Z})$.*

Our construction is given in the next section. Sections 3 and 4 provide some preliminary results and the proof of Theorem 1 respectively.

2. THE CONSTRUCTION

For $x = (x_1, x_2, \dots, x_n) \in [q]^n$, we will write $P(x) = \mathbb{P}(X_1 = x_1, \dots, X_n = x_n)$. To motivate the construction, we begin by noting that the finite-dimensional distributions P of the 4-coloring in [2] are defined recursively by $P(\emptyset) = 1$ and

$$(1) \quad P(x) = \frac{1}{2(n+1)} \sum_{i=1}^n P(\hat{x}_i)$$

Date: July 15, 2014.

for proper $x \in [4]^n$, where \hat{x}_i is obtained from x by deleting the i th entry in x . Of course, even if x is proper, \hat{x}_i may not be. So the definition is completed by setting $P(x) = 0$ for x 's that are not proper.

For general $q \geq 4$, we will now allow the coefficients in the defining sum to depend on i as well as n . Considering many special cases, and the constraints imposed by the 1-dependence requirement, we were led to define

$$(2) \quad P(x) = \frac{1}{D(n+1)} \sum_{i=1}^n C(n-2i+1)P(\hat{x}_i)$$

for proper $x \in [4]^n$, in terms of two sequences C and D . Again motivated by computations in special cases, we take

$$\begin{aligned} C(n) &= T_n(\sqrt{q}/2), & n \geq 0; \\ D(n) &= \sqrt{q}U_{n-1}(\sqrt{q}/2), & n \geq 1, \end{aligned}$$

where T_n and U_n are the Chebyshev polynomials of the first and second kind respectively.

There are several standard equivalent definitions of Chebyshev polynomials. One is

$$(3) \quad T_n(u) = \cosh(nt) \quad \text{and} \quad U_n(u) = \frac{\sinh[(n+1)t]}{\sinh(t)}, \quad \text{where } u = \cosh(t).$$

A variant definition using trigonometric functions (e.g. (22:3:3-4) of [3]) is easily seen to be equivalent by taking t imaginary; the hyperbolic function version is convenient for arguments $u \geq 1$. Another definition is

$$T_n(u) = \sum_{k=0}^{\lfloor \frac{n}{2} \rfloor} \binom{n}{2k} u^{n-2k} (u^2 - 1)^k \quad \text{and} \quad U_n(u) = \sum_{k=0}^{\lfloor \frac{n}{2} \rfloor} \binom{n+1}{2k+1} u^{n-2k} (u^2 - 1)^k.$$

That this is equivalent to (3) follows from e.g. (22:3:1-2) of [3].

If x is not a proper coloring, we take $P(x) = 0$ as before. We extend both sequences C and D to all integer arguments by taking $C(n)$ and $D(n)$ to be even and odd functions of n respectively (in accordance with (3)).

Observe that $C(n)$ and $D(n)$ are strictly positive for $q \geq 4$ and $n \geq 1$, and therefore $P(x)$ is strictly positive for all proper x . Note also that $C(n-2i+1)/D(n+1)$ is rational; therefore so is $P(x)$. (The factors of \sqrt{q} cancel). When $q = 4$ we have $C(n) = 1$ and $D(n) = 2n$, and so (2) reduces to (1) in this case. As we will see, the fact that the coefficients in (2) depend on i substantially complicates the verifications of the required properties of P .

Here are a few examples of cylinder probabilities generated by (2).

$$\begin{aligned} P(1) &= \frac{1}{q}, & P(12) &= \frac{1}{q(q-1)}, & P(121) &= \frac{1}{q^2(q-1)}, & P(123) &= \frac{1}{q^2(q-2)}, \\ P(1212) &= \frac{q-3}{q^2(q-1)(q^2-3q+1)}, & P(1234) &= \frac{1}{q^2(q^2-3q+1)}. \end{aligned}$$

3. PRELIMINARY RESULTS

Chebyshev polynomials satisfy a number of standard identities. They lead to identities satisfied by the sequences C and D . The first three in the proposition below are examples of this. The fourth is a consequence of the third one. Before stating them, we record some values of C and D to facilitate checking computations here and later.

$$C(0) = 1, \quad C(1) = \frac{\sqrt{q}}{2}, \quad C(2) = \frac{q-2}{2}, \quad C(3) = \frac{\sqrt{q}(q-3)}{2}, \quad C(4) = \frac{q^2-4q+2}{2}.$$

$$D(0) = 0, \quad D(1) = \sqrt{q}, \quad D(2) = q, \quad D(3) = \sqrt{q}(q-1), \quad D(4) = q(q-2).$$

Proposition 2. *For $j, k, \ell, m, n \in \mathbb{Z}$, the following identities hold.*

- (4) $2C(m)C(n) = C(m+n) + C(n-m).$
- (5) $\frac{q-4}{2q}D(m)D(n) = C(m+n) - C(n-m).$
- (6) $2C(m)D(n) = D(m+n) + D(n-m).$
- (7) $C(j+k)D(k+\ell) = C(k)D(j+k+\ell) - C(\ell)D(j).$

Proof. The first three parts are immediate consequences of (22:5:5-7) in [3], or 22.7.24-26 in [1], if m and n are nonnegative. None of the identities is changed by changing the sign of either m or n . Therefore, they hold for all m and n . Alternatively, the identities may be checked directly from (3) using the product formulae for hyperbolic functions. For (7), replace the products of C 's and D 's by sums of D 's using (6), and then use the fact that D is an odd function. \square

Next we verify some identities that involve both the sequences C and D and the measure P defined by (2). For the statement of the second part of the next result, let

$$Q(x) = \frac{1}{D(n+1)} \sum_{i=1}^n C(2i)P(\hat{x}_i) \quad \text{and} \quad Q^*(x) = \frac{1}{D(n+1)} \sum_{i=1}^n C(2n-2i+2)P(\hat{x}_i)$$

for $x \in [q]^n$. The first part of Proposition 3 is needed in proving the second part, which plays a key role in the proof of consistency and 1-dependence of P . Note the similarity between the left side of (8) and the right side of (2).

Proposition 3. *If $n \geq 1$, and x is a proper coloring of length n , then*

- (8) $\sum_{i=1}^n D(n-2i+1)P(\hat{x}_i) = 0;$
- (9) $Q(x) = Q^*(x) = P(x)C(n+1).$

Proof. For the first statement, let R be the set of proper colorings, and \hat{x}_A be obtained by deleting the entries x_i for $i \in A$ from x . The proof of (8) is by induction on n , the length of x . The identity is easily seen to be true if $n \leq 2$. Suppose that (8) is true for all x of length $n-1$, and let $x \in R$ have length n . For those i with $\hat{x}_i \in R$, applying (8) gives

$$(10) \quad \sum_{j=1}^{i-1} D(n-2j)P(\hat{x}_{i,j}) + \sum_{j=i+1}^n D(n-2j+2)P(\hat{x}_{i,j}) = 0.$$

On the other hand, if $\hat{x}_i \notin R$, then $1 < i < n$ and

$$(11) \quad P(\hat{x}_{i,j}) = 0 \text{ if } |j - i| > 1 \text{ and } P(\hat{x}_{i-1,i}) = P(\hat{x}_{i,i+1}).$$

The left side of (8) for x can be written, using the definition of $P(\hat{x}_i)$ and then (6), as

$$(12) \quad \begin{aligned} &= \frac{1}{D(n)} \sum_{\substack{1 \leq i \leq n: \\ \hat{x}_i \in R}} D(n - 2i + 1) \left[\sum_{1 \leq j < i} C(n - 2j) P(\hat{x}_{i,j}) + \sum_{i < j \leq n} C(n - 2j + 2) P(\hat{x}_{i,j}) \right] \\ &= \frac{1}{2D(n)} \sum_{\substack{1 \leq j < i \leq n: \\ \hat{x}_i \in R}} [D(2n - 2i - 2j + 1) + D(2j - 2i + 1)] P(\hat{x}_{i,j}) \\ &\quad + \frac{1}{2D(n)} \sum_{\substack{1 \leq i < j \leq n: \\ \hat{x}_i \in R}} [D(2n - 2i - 2j + 3) + D(2j - 2i - 1)] P(\hat{x}_{i,j}). \end{aligned}$$

Rearranging, and ignoring the $2D(n)$ in the denominator, gives

$$(13) \quad \begin{aligned} &\sum_{i=1}^n \mathbf{1}[\hat{x}_i \in R] \left[\sum_{j=1}^{i-1} [D(2n - 2i - 2j + 1) + D(2j - 2i + 1)] P(\hat{x}_{i,j}) \right. \\ &\quad \left. + \sum_{j=i+1}^n [D(2n - 2i - 2j + 3) + D(2j - 2i - 1)] P(\hat{x}_{i,j}) \right]. \end{aligned}$$

We must show that (10) and (11) imply that (13) is zero.

We would like to write (13) as a linear combination of expressions that vanish because of (10) and (11) as follows.

$$(14) \quad \sum_{\substack{1 \leq i \leq n: \\ \hat{x}_i \in R}} \alpha_i \left[\sum_{j=1}^{i-1} D(n - 2j) P(\hat{x}_{i,j}) + \sum_{j=i+1}^n D(n - 2j + 2) P(\hat{x}_{i,j}) \right] + \sum_{\substack{1 \leq i \leq n: \\ \hat{x}_i \in R}} \sum_{j=1}^n \beta_{i,j} P(\hat{x}_{i,j}),$$

where $\beta_{i,i} = \beta_{i,i-1} + \beta_{i,i+1} = 0$. If $1 \leq i < j \leq n$, the coefficient of $P(\hat{x}_{i,j})$ in (13) is

$$(15) \quad \begin{aligned} &\mathbf{1}[\hat{x}_j \in R] [D(2n - 2i - 2j + 1) + D(2i - 2j + 1)] \\ &+ \mathbf{1}[\hat{x}_i \in R] [D(2n - 2i - 2j + 3) + D(2j - 2i - 1)]. \end{aligned}$$

The coefficient of $P(\hat{x}_{i,j})$ in (14) is

$$(16) \quad \mathbf{1}[\hat{x}_j \in R] \alpha_j D(n - 2i) + \mathbf{1}[\hat{x}_i \in R] \alpha_i D(n - 2j + 2) + \mathbf{1}[\hat{x}_i \notin R] \beta_{i,j} + \mathbf{1}[\hat{x}_j \notin R] \beta_{j,i}.$$

We need to choose the α 's and β 's so that (15) and (16) agree. If $\hat{x}_i, \hat{x}_j \in R$, this says

$$D(2n - 2i - 2j + 1) + D(2n - 2i - 2j + 3) = \alpha_j D(n - 2i) + \alpha_i D(n - 2j + 2)$$

since D is an odd function. It may sound unreasonable to expect to solve this system, since there are n unknowns and $\binom{n}{2}$ equations. However, D satisfies relations that make this possible. Solving the equations for small n suggests trying $\alpha_i = 2C(n - 2i + 1)$. The fact that this choice solves these equations for all choices of n, i, j then follows from (6)

and the fact that D is odd. If $\hat{x}_i \notin R$ and $\hat{x}_j \notin R$, (15) and (16) agree if $\beta_{i,j} + \beta_{j,i} = 0$. If $\hat{x}_i \in R$ and $\hat{x}_j \notin R$, they agree if

$$D(2n - 2i - 2j + 3) + D(2j - 2i - 1) = \alpha_i D(n - 2j + 2) + \beta_{j,i}.$$

Using (6) again gives $\beta_{j,i} = 2D(2j - 2i - 1)$. Similarly, if $\hat{x}_i \notin R$ and $\hat{x}_j \in R$, they agree if $\beta_{i,j} = 2D(2i - 2j + 1)$. With these choices, β is anti-symmetric, and $\beta_{k,k-1} = 2D(1)$ and $\beta_{k,k+1} = 2D(-1)$, so $\beta_{k,k-1} + \beta_{k,k+1} = 0$ as required. This completes the induction argument.

For (9), consider the case of Q first. Use the definition of P to write the right side of (9) as

$$\frac{C(n+1)}{D(n+1)} \sum_{i=1}^n C(n-2i+1)P(\hat{x}_i).$$

Using (4), this becomes

$$\frac{1}{2D(n+1)} \sum_{i=1}^n C(2n-2i+2)P(\hat{x}_i) + \frac{1}{2}Q(x).$$

Therefore, we need to prove that

$$\sum_{i=1}^n [C(2n-2i+2) - C(2i)] P(\hat{x}_i) = 0.$$

But by (5), this follows from (8). The proof for Q^* is similar. \square

4. PROOF OF THE MAIN RESULT

We will often write $x_1 x_2 \cdots x_n$ instead of (x_1, x_2, \dots, x_n) below. If $x \in [q]^m$ and $y \in [q]^n$, let xy denote the word $x_1 \cdots x_m y_1 \cdots y_n \in [q]^{m+n}$.

Proof of Theorem 1. We first need to show that the finite dimensional distributions defined in (2) are consistent, i.e., that

$$(17) \quad \sum_{a \in [q]} P(xa) = P(x), \quad x \in [q]^n, \quad n \geq 0.$$

This is true if x is not proper, since then xa is also not proper, and so both sides vanish. For proper x , the proof is by induction on n . Note that for $a \in [q]$,

$$P(a) = \frac{C(0)}{D(2)} = \frac{1}{q},$$

so $\sum_{a \in [q]} P(a) = 1$. This gives (17) for $n = 0$. Suppose it holds for all $x \in [q]^{n-1}$ with $n \geq 1$. Then for proper $x \in [q]^n$, using the induction hypothesis in the second equality,

$$\begin{aligned} \sum_{a \in [q]} P(xa) &= \sum_{a \neq x_n} \frac{1}{D(n+2)} \left[\sum_{i=1}^n C(n-2i+2)P(\hat{x}_i a) + C(-n)P(x) \right] \\ &= \frac{1}{D(n+2)} \left[\sum_{i=1}^n C(n-2i+2)P(\hat{x}_i) - C(-n+2)P(x) + (q-1)C(-n)P(x) \right]. \end{aligned}$$

The middle term in the second line accounts for the missing term $a = x_n$ when the inductive hypothesis is applied to the case $i = n$ (since $\widehat{x}_n x_n = x$). Using $(j, k, \ell) = (1, n - 2i + 1, 2i)$ in (7) gives

$$\frac{C(n - 2i + 2)}{D(n + 2)} = \frac{C(n - 2i + 1)}{D(n + 1)} - \frac{C(2i)D(1)}{D(n + 2)D(n + 1)}.$$

Therefore

$$\sum_{a \in [q]} P(xa) = P(x) - \frac{Q(x)}{D(n + 2)} - \frac{C(n - 2)}{D(n + 2)}P(x) + (q - 1)\frac{C(n)}{D(n + 2)}P(x).$$

This is $P(x)$, as required, by (9) and the fact that

$$(q - 1)C(n) = C(n - 2) + C(n + 1)D(1),$$

which is obtained by taking $(j, k, \ell) = (2, -n, n + 1)$ in (7), and then canceling a factor of \sqrt{q} .

Invariance of the measure under permutations of colors and translations is immediate from the definition. Invariance under reflection amounts to checking $P(x) = P(x_n \cdots x_1)$, which follows from the fact that the coefficients of \widehat{x}_i and \widehat{x}_{n-i+1} in (2), which are $C(n - 2i + 1)$ and $C(-n + 2i - 1)$ respectively, are equal by the symmetry of C .

For 1-dependence, we need to show that for $x \in [q]^m$ and $y \in [q]^n$ with $m, n \geq 0$,

$$P(x * y) = P(x)P(y),$$

where the $*$ means that there is no constraint at the single site between x and y . This is again true if x or y is not proper since then both sides are zero. For proper x and y , the proof is by induction, but now on $m + n$. The statement is immediate if $m = 0$ or $n = 0$. So, we take $m \geq 1$ and $n \geq 1$.

There are two cases, according to whether or not xy is a proper coloring, i.e., whether x_m and y_1 are equal or different. Assume first that $x_m = y_1$. Without loss of generality, take their common value to be 1. Then using the definition of P , including the fact that $P(xy) = 0$,

$$\begin{aligned} (18) \quad P(x * y) &= \sum_{a \in [q]} P(xay) = \frac{1}{D(n + m + 2)} \sum_{a \neq 1} \left[\sum_{i=1}^m C(n + m - 2i + 2)P(\widehat{x}_i a y) \right. \\ &\quad \left. + C(n - m)P(xy) + \sum_{j=1}^n C(n - m - 2j)P(x a \widehat{y}_j) \right] \\ &= \frac{1}{D(n + m + 2)} \left[\sum_{i=1}^m C(n + m - 2i + 2)P(\widehat{x}_i * y) + \sum_{j=1}^n C(n - m - 2j)P(x * \widehat{y}_j) \right]. \end{aligned}$$

Using the induction hypothesis, this becomes

$$P(x * y) = \frac{1}{D(n + m + 2)} \left[P(y) \sum_{i=1}^m C(n + m - 2i + 2)P(\widehat{x}_i) + P(x) \sum_{j=1}^n C(n - m - 2j)P(\widehat{y}_j) \right].$$

Taking $(j, k, l) = (n, m - 2i + 1, i)$ in (7) gives

$$\frac{C(n + m - 2i + 2)}{D(n + m + 2)} = \frac{C(m - 2i + 1)}{D(m + 1)} - \frac{C(2i)D(n + 1)}{D(m + 1)D(n + m + 2)}.$$

Similarly,

$$\frac{C(m + 2j - n)}{D(n + m + 2)} = \frac{C(2j - n - 1)}{D(n + 1)} - \frac{C(2n - 2j + 2)D(m + 1)}{D(n + 1)D(n + m + 2)}.$$

Therefore, since $C(\cdot)$ is even,

$$P(x * y) = P(y) \left[P(x) - \frac{D(n + 1)}{D(n + m + 2)} Q(x) \right] + P(x) \left[P(y) - \frac{D(m + 1)}{D(n + m + 2)} Q^*(y) \right].$$

By (9),

$$P(x * y) = P(x)P(y) \left[2 - \frac{C(m + 1)D(n + 1) + C(n + 1)D(m + 1)}{D(n + m + 2)} \right].$$

Taking $(j, k, l) = (n, m - 2i + 1, i)$ in (7), we see that the expression in brackets above is 1, as required.

Assume now that $x_m \neq y_1$, say $x_m = 1$ and $y_1 = 2$. Then

$$(19) \quad P(x * y) = \sum_{a \in [q]} P(xay) = \frac{1}{D(n + m + 2)} \sum_{a \neq 1, 2} \left[\sum_{i=1}^m C(n + m - 2i + 2) P(\widehat{x}_i ay) \right. \\ \left. + C(n - m) P(xy) + \sum_{j=1}^n C(n - m - 2j) P(xa\widehat{y}_j) \right] \\ \frac{1}{D(n + m + 2)} \left[\sum_{i=1}^m C(n + m - 2i + 2) P(\widehat{x}_i * y) + \sum_{j=1}^n C(n - m - 2j) P(x * \widehat{y}_j) \right]$$

as in the previous case. However, in the previous case, the term $P(xy)$ dropped out because xy was not a proper coloring. In this case, the term $(q - 2)C(n - m)P(xy)$ is cancelled by the terms $-P(xy)C(n - m + 2)$ and $-P(xy)C(n - m - 2)$, which arise from

$$\sum_{a \neq 1, 2} P(\widehat{x}_m ay) = P(\widehat{x}_m * y) - P(xy) \quad \text{and} \quad \sum_{a \neq 1, 2} P(xa\widehat{y}_1) = P(x * \widehat{y}_1) - P(xy).$$

The fact that the overall coefficient of $P(xy)$ vanishes is a consequence (4) with $m = 2$, since $2C(2) = q - 2$. The rest of the proof is the same as in the case $x_m = y_1$ above. \square

REFERENCES

- [1] Abramowitz, M. and Stegun, I.A., editors. Handbook of Mathematical Functions with Formulas, Graphs and Mathematical Tables. National Bureau of Standards, 1972.
- [2] Holroyd, A. E. and Liggett, T. M. Finitely dependent coloring. arXiv:1403.2448
- [3] Oldham, K., Myland, J. and Spanier, J. *An Atlas of Functions*. Second edition. Springer, New York, 2009.

(Alexander E. Holroyd) MICROSOFT RESEARCH

(Thomas M. Liggett) UNIVERSITY OF CALIFORNIA, LOS ANGELES